

# Probabilistic Model Checking and Strategy Synthesis

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# Probabilistic Model Checking and Strategy Synthesis

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Joint work with: Marta Kwiatkowska, Vojtěch Forejt, Gethin Norman, Hongyang Qu, Aistis Simaitis, Taolue Chen, Bruno Lacerda, Nick Hawes

### Overview

#### Probabilistic model checking

- verification vs. strategy synthesis
- Markov decision processes (MDPs)
- example: robot navigation

#### Multi-objective probabilistic model checking

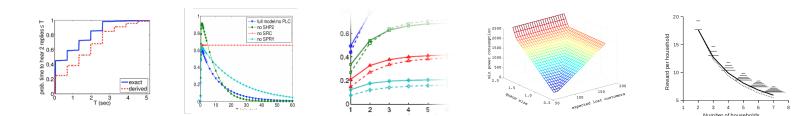
- examples: power management/team-formation
- Stochastic (multi-player) games
  - example: energy management

## Motivation

- Verifying probabilistic systems...
  - unreliable or unpredictable behaviour
    - · failures of physical components
    - message loss in wireless communication
    - unreliable sensors/actuators
  - randomisation in algorithms/protocols
    - random back-off in communication protocols
    - random routing to reduce flooding or provide anonymity
- We need to verify quantitative system properties
  - "the probability of the airbag failing to deploy within 0.02 seconds of being triggered is at most 0.001"
  - not just correctness: reliability, timeliness, performance, ...
  - not just verification: correctness by construction

# Probabilistic model checking

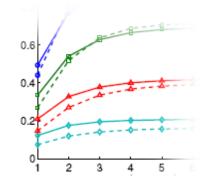
- Construction and analysis of probabilistic models
  - state-transition systems labelled with probabilities (e.g. Markov chains, Markov decision processes)
  - from a description in a high-level modelling language
- Properties expressed in temporal logic, e.g. PCTL:
  - trigger  $\rightarrow$  P<sub> $\geq 0.999$ </sub> [ F<sup> $\leq 20$ </sup> deploy ]
  - "the probability of the airbag deploying within 20ms of being triggered is at at least 0.999"
  - properties checked against models using exhaustive search and numerical computation



0.5 \$0.4

# Probabilistic model checking

- Many types of probabilistic models supported
- Wide range of quantitative properties, expressible in temporal logic (probabilities, timing, costs, rewards, ...)
- Often focus on numerical results (probabilities etc.)
  - analyse trends, look for system flaws, anomalies
    - P<sub>≤0.1</sub> [F *fail*] "the probability of a failure occurring is at most 0.1"

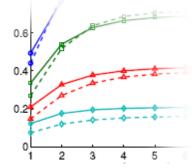


 P<sub>=?</sub> [F fail] – "what is the probability of a failure occurring?"

# Probabilistic model checking

- Many types of probabilistic models supported
- Wide range of quantitative properties, expressible in temporal logic (probabilities, timing, costs, rewards, ...)
- Often focus on numerical results (probabilities etc.)
   analyse trends, look for system flaws, anomalies
- Provides "exact" numerical results/guarantees
  - compared to, for example, simulation
- Combines numerical & exhaustive analysis

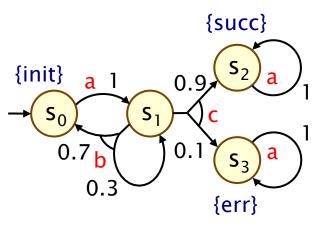
   especially useful for nondeterministic models



- Fully automated, tools available, widely applicable
  - network/communication protocols, security, biology, robotics & planning, power management, ...

## Markov decision processes (MDPs)

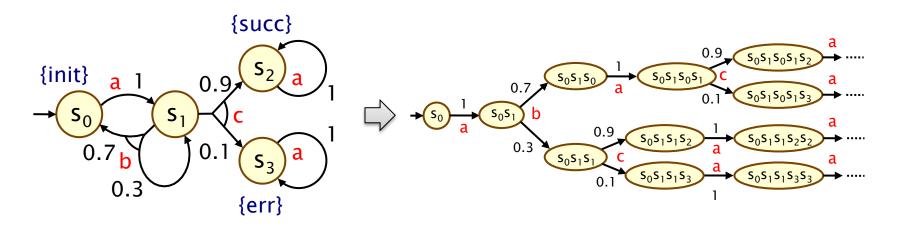
- Markov decision processes (MDPs)
  - widely used also in: AI, planning, optimal control, ...
  - model nondeterministic as well as probabilistic behaviour



- Nondeterminism for:
  - control: decisions made by a controller or scheduler
  - adversarial behaviour of the environment
  - concurrency/scheduling: interleavings of parallel components
  - abstraction, or under-specification, of unknown behaviour

# Strategies

- A strategy (or "policy" or "adversary")
  - is a resolution of nondeterminism, based on history
  - i.e. a mapping from finite paths to (distributions over) actions
  - induces (infinite-state) Markov chain (and probability space)

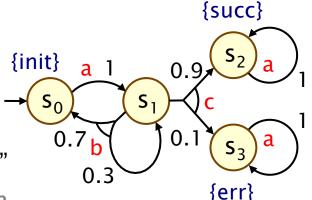


- Classes of strategies:
  - memory: memoryless, finite-memory, or infinite-memory
  - randomisation: deterministic or randomised

# Verification vs. Strategy synthesis

#### • 1. Verification

- quantify over all possible strategies (i.e. best/worst-case)
- $P_{\leq 0.1}$  [F *err*]: "the probability of an error occurring is  $\leq 0.1$  for all strategies"



 applications: randomised communication protocols, randomised distributed algorithms, security, ...

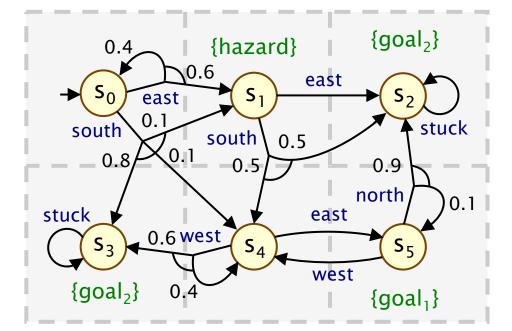
#### 2. Strategy synthesis

- generation of "correct-by-construction" controllers
- $P_{\leq 0.1}$  [F *err*]: "does there exist a strategy for which the probability of an error occurring is  $\leq 0.1$ ?"
- applications: robotics, power management, security, ...
- Two dual problems; same underlying computation:
  - compute optimal (minimum or maximum) values

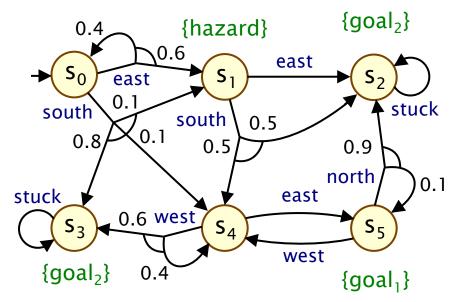
### Running example

#### • Example MDP

- robot moving through terrain divided in to 3 x 2 grid



## Example – Reachability



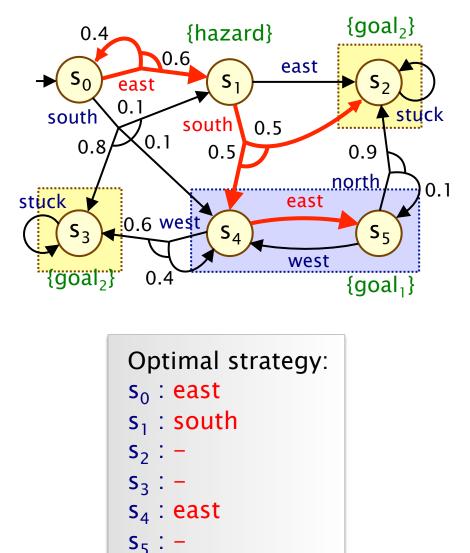
Verify:  $P_{\leq 0.6}$  [ F goal<sub>1</sub> ] or Synthesise for:  $P_{\geq 0.4}$  [ F goal<sub>1</sub> ]  $\Downarrow$ Compute:  $P_{max=?}$  [ F goal<sub>1</sub> ]

Optimal strategies: memoryless and deterministic

Computation:

graph analysis + numerical soln. (linear programming, value iteration, policy iteration)

## Example – Reachability



Verify:  $P_{\leq 0.6}$  [F goal<sub>1</sub>] or Synthesise for:  $P_{\geq 0.4}$  [F goal<sub>1</sub>]  $\Downarrow$ Compute:  $P_{max=?}$  [F goal<sub>1</sub>] = 0.5

Optimal strategies: memoryless and deterministic

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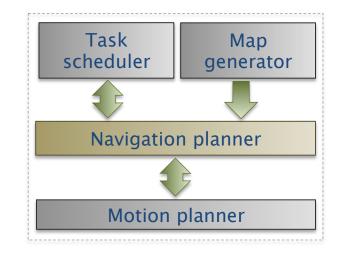
## MDPs – Other properties

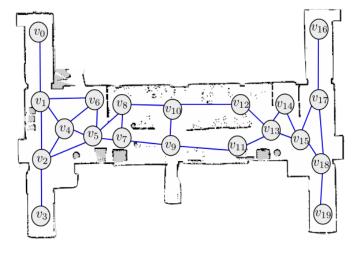
- Costs and rewards (expected, accumulated values)
  - e.g. R<sub>min=?</sub> [F goal<sub>2</sub>] "what is the minimum expected number of moves needed to reach goal<sub>2</sub>?"
  - optimal strategies: memoryless and deterministic
  - similar computation to probabilistic reachability
- Probabilistic LTL (multiple temporal operators)
  - e.g.  $P_{max=?}$  [ (G¬hazard)  $\land$  (GF goal<sub>1</sub>) ] "maximum probability of avoiding hazard and visiting goal<sub>1</sub> infinitely often?"
  - optimal strategies: finite-memory and deterministic
  - build product MDP, graph analysis, probabilistic reachability
- Expected cost/reward to satisfy (co-safe) LTL formula
  - e.g.  $R_{min=?}$  [ $\neg zone_3 U (zone_1 \land (F zone_4))$ ] "minimise exp. time to patrol zones 1 then 4, without passing through 3".

# Application: Robot navigation

- Navigation planning: [IROS'14]
  - MDP models navigation through an uncertain environment
  - LTL used to formally specify tasks to be executed
  - synthesise finite-memory strategies to construct plans/controllers
  - links to continuous-space planner





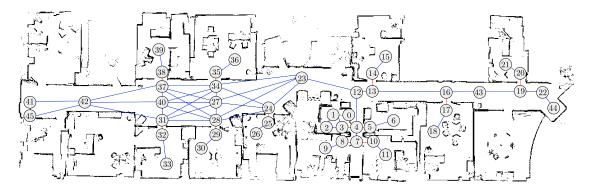


# Application: Robot navigation

- Navigation planning MDPs
  - expected timed on edges + probabilities
  - learnt using data from previous explorations
- LTL-based task specification



- expected time to satisfy (one or more) co-safe LTL formulas
- c.f. ad-hoc reward structures, e.g. with discounting
- also: efficient re-planning [IROS'14]; progress metric [IJCAI'15]
- Implementation
  - MetraLabs Scitos A5 robot + ROS module based on PRISM



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#### Multi-objective probabilistic model checking

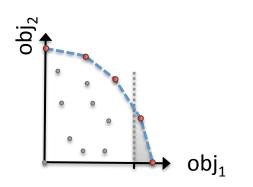
- examples: power management/team-formation
- Stochastic (multi-player) games
  - example: energy management

### Multi-objective model checking

- Multi-objective probabilistic model checking
  - investigate trade-offs between conflicting objectives
  - in PRISM, objectives are probabilistic LTL or expected rewards
- Achievability queries: multi(P<sub>>0.95</sub> [ F send ], R<sup>time</sup><sub>>10</sub> [ C ])
  - e.g. "is there a strategy such that the probability of message transmission is > 0.95 and expected battery life > 10 hrs?"
- Numerical queries: multi(P<sub>max=?</sub> [ F send ], R<sup>time</sup><sub>>10</sub> [ C ])
  - e.g. "maximum probability of message transmission, assuming expected battery life-time is > 10 hrs?"

#### Pareto queries:

- multi(P<sub>max=?</sub>[F send], R<sup>time</sup>max=?[C])
- e.g. "Pareto curve for maximising probability of transmission and expected battery life-time"



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- e.g. "Pareto curve for maximising probability of transmission and expected battery life-time"

obj₁

## Multi-objective model checking

#### Optimal strategies:

- usually finite-memory (e.g. when using LTL formulae)
- may also need to be randomised

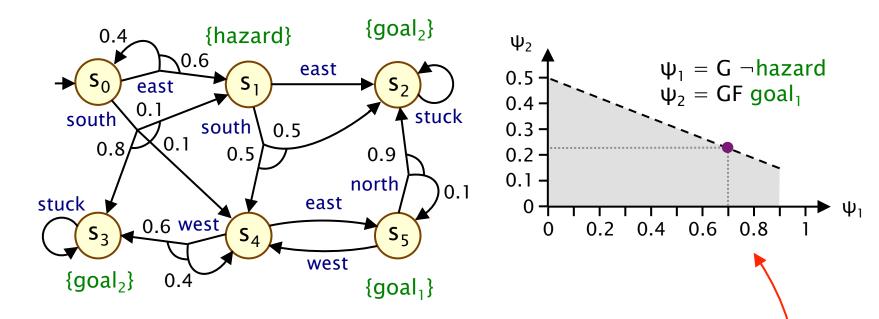
#### Computation:

- construct a product MDP (with several automata), then reduces to linear programming [TACAS'07,TACAS'11]
- can be approximated using iterative numerical methods, via approximation of the Pareto curve [ATVA'12]

#### • Extensions [ATVA'12]

- arbitrary Boolean combinations of objectives
  - · e.g.  $\psi_1 \Rightarrow \psi_2$  (all strategies satisfying  $\psi_1$  also satisfy  $\psi_2$ )
  - (e.g. for assume-guarantee reasoning)
- time-bounded (finite-horizon) properties

#### Example - Multi-objective

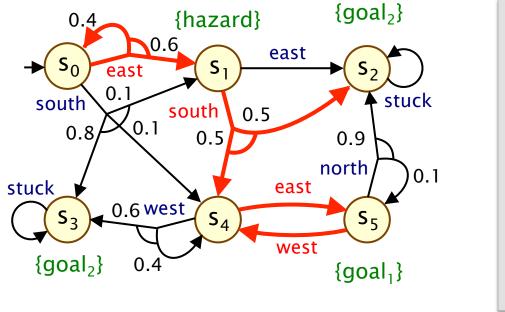


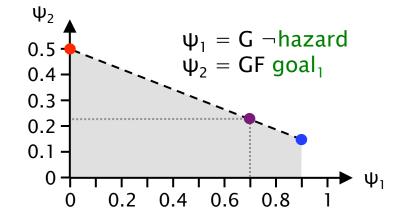
- Achievability query
  - $P_{\geq 0.7}$  [ G ¬hazard ]  $\land$   $P_{\geq 0.2}$  [ GF goal<sub>1</sub> ] ? True (achievable)
- Numerical query

-  $P_{max=?}$  [ GF goal<sub>1</sub> ] such that  $P_{\geq 0.7}$  [ G  $\neg$  hazard ]? ~0.2278

- Pareto query
  - for  $P_{max=?}$  [G ¬hazard]  $\land$   $P_{max=?}$  [GF goal<sub>1</sub>]?

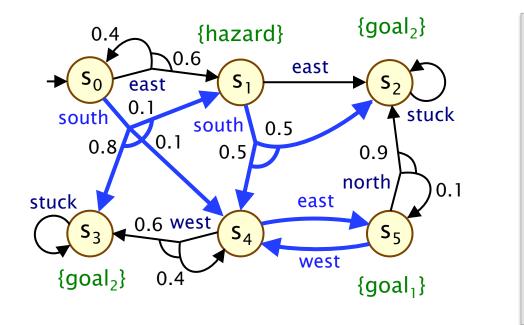
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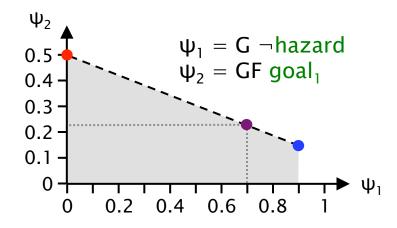




Strategy 1 (deterministic)  $s_0$  : east  $s_1$  : south  $s_2$  :  $s_3$  :  $s_4$  : east  $s_5$  : west

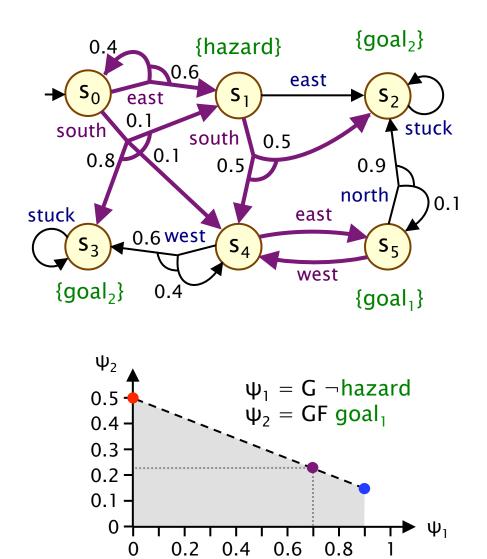
#### Example – Multi–objective





Strategy 2 (deterministic)  $s_0$  : south  $s_1$  : south  $s_2$  :  $s_3$  :  $s_4$  : east  $s_5$  : west

#### Example – Multi–objective



Optimal strategy: (randomised)  $s_0 : 0.3226 : east$  0.6774 : south  $s_1 : 1.0 : south$   $s_2 :$   $s_3 :$   $s_4 : 1.0 : east$  $s_5 : 1.0 : west$ 

## Multi-objective: Applications

#### Synthesis of controllers for dynamic power management [TACAS'11]

#### IBM TravelStar VP disk drive

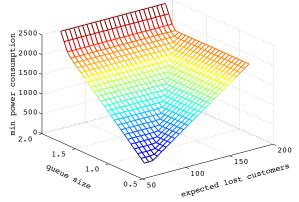
- switches between power modes:
- active/idle/idlelp/stby/sleep

#### MDP model in PRISM<sup>•</sup>

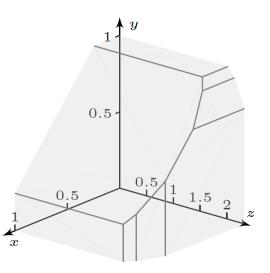
- power manager
- disk requests •
- request queue
- power usage

#### Multi-objective:

"minimise energy consumption, subject to constraints on: (i) expected job queue size; (ii) expected number of lost jobs



#### Synthesis of team formation strategies [CLIMA'11, ATVA'12]



#### Pareto curve:

x="probability of completing task 1"; y="probability of completing task 2"; z="expected size of successful team" 25

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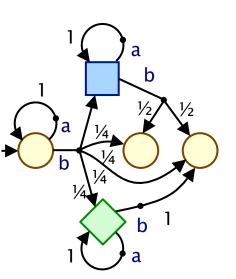
– examples: power management/team-formation

#### Stochastic (multi-player) games

– example: energy management

# Stochastic multi-player games (SMGs)

- Stochastic multi-player games
  - players control states; choose actions
  - models competitive/collaborative behaviour
  - applications: security (system vs. attacker), controller synthesis (controller vs. environment), distributed algorithms/protocols, ...

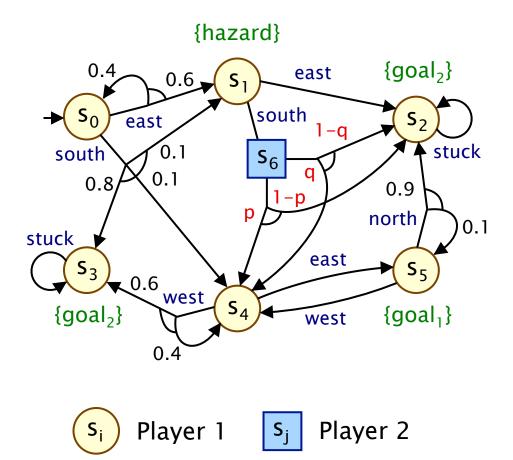


#### Property specifications: rPATL

- $\langle\langle\{1,2\}\rangle\rangle P_{\geq 0.95}$  [  $F^{\leq 45}$  *done* ] : "can nodes 1,2 collaborate so that the probability of the protocol terminating within 45 seconds is at least 0.95, whatever nodes 3,4 do?"
- formally:  $\langle\langle C \rangle\rangle \psi$ : do there exist strategies for players in C such that, for all strategies of other players, property  $\psi$  holds?
- Model checking [TACAS'12,FMSD'13]
  - zero sum properties: analysis reduces to 2-player games
  - PRISM-games: <u>www.prismmodelchecker.org/games</u>

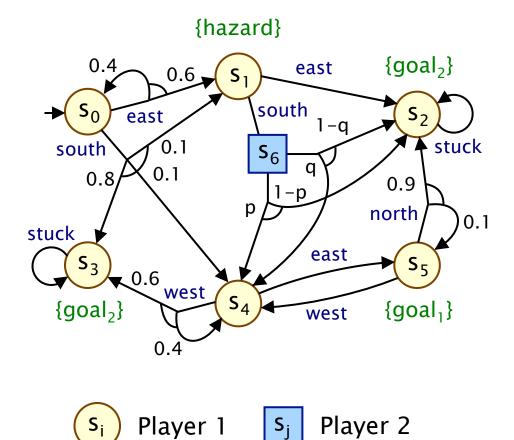
#### Example – Stochastic games

- Two players: 1 (robot controller), 2 (environment)
  - − probability of  $s_1$ −south→ $s_4$  is in [p,q] = [0.5-Δ, 0.5+Δ]



#### Example – Stochastic games

- Two players: 1 (robot controller), 2 (environment)
  - probability of  $s_1$ -south $\rightarrow s_4$  is in [p,q] = [0.5- $\Delta$ , 0.5+ $\Delta$ ]



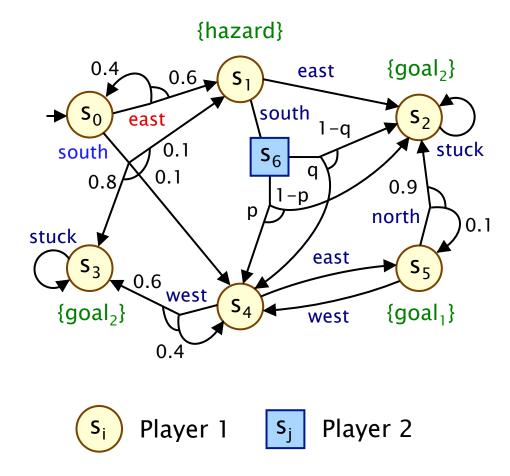
rPATL:  $\langle\langle \{1\}\rangle\rangle P_{max=?} [Fgoal_1]$ 

Optimal strategies: memoryless and deterministic

Computation: graph analysis & numerical approximation

#### Example – Stochastic games

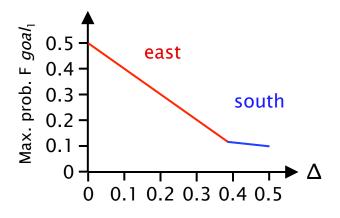
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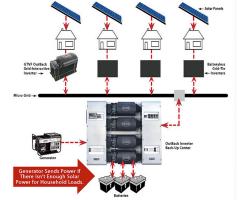


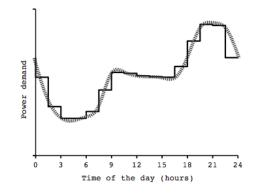
# Application: Energy management

- Energy management protocol for Microgrid
  - randomised demand management protocol
  - random back-off when demand is high
- Original analysis [Hildmann/Saffre'11]
  - protocol increases "value" for clients
  - simulation-based, clients are honest

#### Our analysis

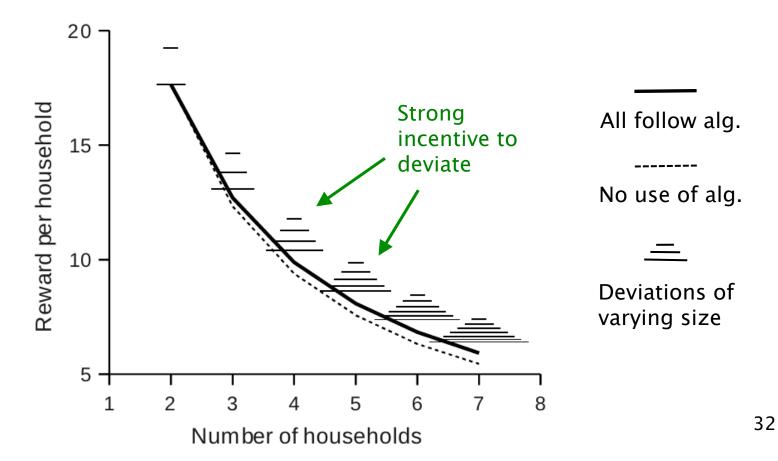
- stochastic multi-player game model
- clients can cheat (and cooperate)
- model checking: PRISM-games
- exposes protocol weakness (incentive for clients to act selfishly
- propose/verify simple fix using penalties





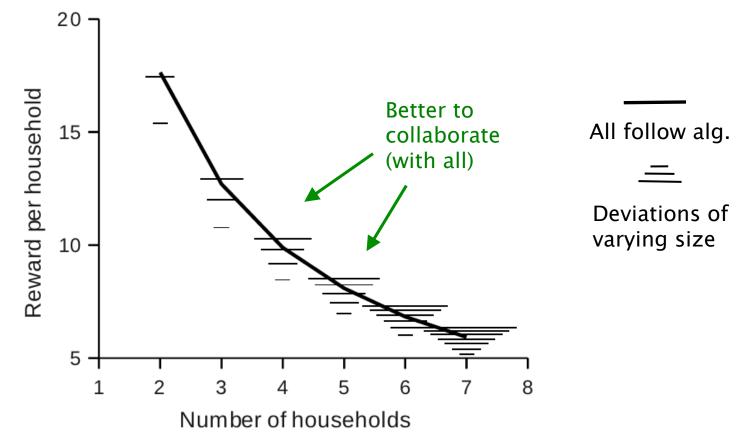
### Results: Competitive behaviour

- Expected total value V per household
  - in rPATL:  $\langle \langle C \rangle \rangle R^{r}C_{max=?}$  [F<sup>0</sup> time=max time] / |C|
  - where  $r_c$  is combined rewards for coalition C



## Results: Competitive behaviour

- Algorithm fix: simple punishment mechanism
  - distribution manager can cancel some loads exceeding  $c_{lim}$



# Conclusion

- Probabilistic model checking
  - verification vs. strategy synthesis
  - Markov decision processes, temporal logic, PRISM

#### Recent directions and extensions

- multi-objective probabilistic model checking
- model checking for stochastic games

#### - Challenges

- stochastic games: multi-objective, equilibria, richer logics
- partial information/observability
- probabilistic models with continuous time (or space)
- scalability, e.g. symbolic methods, abstraction